

# 1 System T

## 1.1 Syntax

$bool$  : **type**.  
 $unary$  : **type**.  
 $term$  : **type**.  
 $clos$  : **type**.  
 $env$  : **type**.  
 $stack$  : **type**.  
 $state$  : **type**.

### Unary number $n$

$0$  :  $unary$ .  
 $n + 1$  :  $unary$ .

### Boolean $b$

$true$  :  $bool$ .  
 $false$  :  $bool$ .

### Term $t$

$n$  :  $term$ .  
 $t_1 t_2$  :  $term$ .  
 $\lambda n. t$  :  $term$ .

### Closure $c$

$(t, \mathcal{E})$  :  $clos$ .

### Environment $\mathcal{E}$

$\square$  :  $env$ .  
 $(\mathcal{E}, n \leftarrow c)$  :  $env$ .

### Stack $\mathcal{S}$

$\square$  :  $stack$ .  
 $c : \mathcal{S}$  :  $stack$ .

### State $\sigma$

$\langle t, \mathcal{E}, \mathcal{S} \rangle$  :  $state$ .

## 1.2 Judgments

$n_1 = n_2 \rightsquigarrow b$  : **type**.  
 $\mathcal{E}(n) = c$  : **type**.  
 $\mathcal{E}(n_1)_{n_2} = c$  : **type**.  
 $\sigma_1 \rightarrow \sigma_2$  : **type**.  
 $b?c_1:c_2 \rightsquigarrow c$  : **type**.

### 1.3 Equality

$$\frac{}{0 = 0 \rightsquigarrow \mathbf{true}} [\text{Equal}_1]$$

$$\frac{}{0 = n + 1 \rightsquigarrow \mathbf{false}} [\text{Equal}_2]$$

$$\frac{}{n + 1 = 0 \rightsquigarrow \mathbf{false}} [\text{Equal}_3]$$

$$\frac{n = m \rightsquigarrow b}{n + 1 = m + 1 \rightsquigarrow b} [\text{Equal}_4]$$

**%mode**  $+n = +m \rightsquigarrow -b$   
**%worlds**  $() \quad n = m \rightsquigarrow b$   
**%terminates**  $n \quad n = m \rightsquigarrow b$   
**%unique**  $+n = +m \rightsquigarrow -1b$

### 1.4 Conditional

$$\frac{}{\mathbf{true}?c_1:c_2 \rightsquigarrow c_1} [\text{Cond\_1}]$$

$$\frac{}{\mathbf{false}?c_1:c_2 \rightsquigarrow c_2} [\text{Cond\_2}]$$

**%mode**  $+b?+c:+c' \rightsquigarrow -c''$   
**%worlds**  $() \quad b?c:c' \rightsquigarrow c''$   
**%terminates**  $b \quad b?c_1:c_2 \rightsquigarrow c'$   
**%unique**  $+b?+c:+c' \rightsquigarrow -1c''$

### 1.5 Fetch

$$\frac{n = m \rightsquigarrow b \quad \mathcal{E}(m) = c' \quad b?c:c' \rightsquigarrow c''}{(\mathcal{E}, n \leftarrow c)(m) = c''} [\text{Fetch\_1}]$$

**%mode**  $+\mathcal{E}(+n) = -c$   
**%worlds**  $() \quad \mathcal{E}(n) = c$   
**%terminates**  $\mathcal{E} \quad \mathcal{E}(n) = c$   
**%unique**  $+\mathcal{E}(+n) = -1c$

### 1.6 Evaluation

$$\frac{\mathcal{E}(x) = (t, \mathcal{E}')}{\langle x, \mathcal{E}, \mathcal{S} \rangle \rightarrow \langle t, \mathcal{E}', \mathcal{S} \rangle} [\text{Eval\_Var}]$$

$$\frac{}{\langle (t_1 t_2), \mathcal{E}, \mathcal{S} \rangle \rightarrow \langle t_1, \mathcal{E}, (t_2, \mathcal{E}) : \mathcal{S} \rangle} [\text{Eval\_App}]$$

$$\frac{}{\langle \lambda x.t, \mathcal{E}, c : \mathcal{S} \rangle \rightarrow \langle t, (\mathcal{E}, x \leftarrow c), \mathcal{S} \rangle} [\text{Eval\_Abs}]$$

**%mode**  $+\sigma_1 \rightarrow -\sigma_2$   
**%worlds**  $() \quad \sigma_1 \rightarrow \sigma_2$   
**%unique**  $+\sigma_1 \rightarrow -1\sigma_2$