

1 Simply typed λ -calculus

1.1 Syntax

1.1.1 Types

$type$: **type**.
 \mathbf{unit} : $type$.
 $\sqcup \rightarrow \sqcup$: $type \rightarrow type \rightarrow type$.
 $\sqcup \times \sqcup$: $type \rightarrow type \rightarrow type$.

1.1.2 Terms

$term$: **type**.
 $\langle \rangle$: $term$.
 $\sqcup \sqcup$: $term \rightarrow term \rightarrow term$.
 $\lambda_{\sqcup:\sqcup.\sqcup}$: $type \rightarrow (term \rightarrow term) \rightarrow term$.
%binding $1 \mapsto 3$ **in** $\lambda_{\sqcup:\sqcup.\sqcup}$
 $\langle \sqcup, \sqcup \rangle$: $term \rightarrow term \rightarrow term$.
let $\langle \sqcup, \sqcup \rangle = \sqcup$ **in** \sqcup : $term \rightarrow (term \rightarrow term \rightarrow term) \rightarrow term$.
%binding $2 \mapsto 4$ **in** **let** $\langle \sqcup, \sqcup \rangle = \sqcup$ **in** \sqcup
%binding $1 \mapsto 4$ **in** **let** $\langle \sqcup, \sqcup \rangle = \sqcup$ **in** \sqcup

1.2 Typing judgment

$\vdash_{\sqcup:\sqcup}$: $term \rightarrow type \rightarrow \mathbf{type}$.

$$\overline{\vdash \langle \rangle : \mathbf{unit}}^{[\text{of_empty}]}$$

$$\frac{\{x\} \vdash x : \tau \rightarrow \vdash t[x] : \tau'}{\vdash \lambda x : \tau. t[x] : \tau \rightarrow \tau'}^{[\text{of_lam}]}$$

$$\frac{\vdash t_1 : \tau \rightarrow \tau' \quad \vdash t_2 : \tau}{\vdash t_1 t_2 : \tau'} [\text{of_app}]$$

$$\frac{\vdash t_1 : \tau \quad \vdash t_2 : \tau'}{\vdash \langle t_1, t_2 \rangle : \tau \times \tau'} [\text{of_pair}]$$

$$\frac{\{x\} \vdash x : \tau_1 \rightarrow (\{y\} \vdash y : \tau_2 \rightarrow \vdash u[x][y] : \tau) \quad \vdash t : \tau_1 \times \tau_2}{\vdash \text{let } \langle x, y \rangle = t \text{ in } u[x][y] : \tau} [\text{of_let}]$$

1.2.1 Derived meta-rule for the let macro

$$(\{x\} \vdash x : \tau \rightarrow \vdash t[x] : \tau') \quad \wedge \quad \vdash t_2 : \tau \quad \Longrightarrow \quad \vdash \lambda x : \tau. t[x] t_2 : \tau' \quad : \quad \text{type.}$$

$$\frac{\frac{\mathcal{D}_{\text{of}_1} \quad \{x\} \vdash x : \tau \rightarrow \vdash t[x] : \tau'}{\vdash \lambda x : \tau. t[x] : \tau \rightarrow \tau'} [\text{of_lam}] \quad \mathcal{D}_{\text{of}_2} \quad \vdash t_2 : \tau}{\vdash \lambda x : \tau. t[x] t_2 : \tau'} [\text{of_app}] \quad [\&]$$

$$\begin{aligned} \% \text{mode} \quad & +\mathcal{D}_{\text{of}_1} \wedge +\mathcal{D}_{\text{of}_2} \Longrightarrow -\mathcal{D}_{\text{of}_3} \\ \% \text{worlds} \quad & () \mathcal{D}_{\text{of}_1} \wedge \mathcal{D}_{\text{of}_2} \Longrightarrow \mathcal{D}_{\text{of}_3} \\ \% \text{total} \quad & \{\} \mathcal{D}_{\text{of}_1} \wedge \mathcal{D}_{\text{of}_2} \Longrightarrow \mathcal{D}_{\text{of}_3} \end{aligned}$$

1.3 Values

$$\sqcup \text{ value} : \text{term} \rightarrow \text{type.}$$

$$\langle \rangle \text{ value} [\text{value_empty}]$$

$$\lambda x : \tau. t[x] \text{ value} [\text{value_lam}]$$

$$\frac{t_1 \text{ value} \quad t_2 \text{ value}}{\langle t_1, t_2 \rangle \text{ value}} [\text{value_pair}]$$

1.4 Reduction semantics

$\sqcup \rightarrow \sqcup \quad : \quad \text{term} \rightarrow \text{term} \rightarrow \text{type}.$

$$\frac{t_1 \rightarrow t'_1}{t_1 t_2 \rightarrow t'_1 t_2} [\text{step_app1}]$$

$$\frac{t_1 \text{ value} \quad t_2 \rightarrow t'_2}{t_1 t_2 \rightarrow t_1 t'_2} [\text{step_app2}]$$

$$\frac{t_1 \rightarrow t'_1}{\langle t_1, t_2 \rangle \rightarrow \langle t'_1, t_2 \rangle} [\text{step_pair1}]$$

$$\frac{t_1 \text{ value} \quad t_2 \rightarrow t'_2}{\langle t_1, t_2 \rangle \rightarrow \langle t_1, t'_2 \rangle} [\text{step_pair2}]$$

$$\frac{t_2 \text{ value}}{\lambda x: \tau. t[x] t_2 \rightarrow t[t_2]} [\text{step_beta}]$$

$$\frac{t_1 \text{ value} \quad t_2 \text{ value}}{\text{let } \langle x, y \rangle = \langle t_1, t_2 \rangle \text{ in } u[x][y] \rightarrow u[t_1][t_2]} [\text{step_let}]$$

1.5 Preservation theorem

$t \rightarrow t' \quad \wedge \quad \vdash t: \tau \quad \implies \quad \vdash t': \tau \quad : \quad \text{type}.$

$$\begin{array}{c}
\frac{\mathcal{D}_{\text{stepE}_1} \quad \mathcal{D}_{\text{ofE}_1} \quad \mathcal{D}_{\text{ofE}'_1}}{t_1 \rightarrow t'_1 \quad \vdash t_1: \tau_2 \rightarrow \tau \quad \vdash t'_1: \tau_2 \rightarrow \tau} \Rightarrow \quad \text{---} [\&_1] \\
\frac{\mathcal{D}_{\text{stepE}_1} \quad \mathcal{D}_{\text{ofE}_1} \quad \mathcal{D}_{\text{ofE}_2}}{t_1 \rightarrow t'_1 \quad \vdash t_1: \tau_2 \rightarrow \tau \quad \vdash t_2: \tau_2} [\text{step_app1}] \quad \wedge \quad \frac{\mathcal{D}_{\text{ofE}_1} \quad \mathcal{D}_{\text{ofE}_2}}{\vdash t_1 t_2: \tau} [\text{of_app}] \quad \Rightarrow \quad \frac{\mathcal{D}_{\text{ofE}'_1} \quad \mathcal{D}_{\text{ofE}_2}}{\vdash t'_1 t_2: \tau} [\text{of_app}]
\end{array}$$

$$\begin{array}{c}
\frac{\mathcal{D}_{\text{stepE}_2} \quad \mathcal{D}_{\text{ofE}_2} \quad \mathcal{D}_{\text{ofE}'_2}}{t_2 \rightarrow t'_2 \quad \vdash t_2: \tau_2 \quad \vdash t'_2: \tau_2} \Rightarrow \quad \text{---} [\&_2] \\
\frac{\mathcal{D}_{\text{valE}_1} \quad \mathcal{D}_{\text{stepE}_2}}{t_1 \text{ value} \quad t_2 \rightarrow t'_2} [\text{step_app2}] \quad \wedge \quad \frac{\mathcal{D}_{\text{ofE}_1} \quad \mathcal{D}_{\text{ofE}_2}}{\vdash t_1: \tau_2 \rightarrow \tau \quad \vdash t_2: \tau_2} [\text{of_app}] \quad \Rightarrow \quad \frac{\mathcal{D}_{\text{ofE}_1} \quad \mathcal{D}_{\text{ofE}'_2}}{\vdash t_1 t_2 \rightarrow t_1 t'_2} [\text{of_app}]
\end{array}$$

$$\begin{array}{c}
\frac{\mathcal{D}_{\text{stepE}_1} \quad \mathcal{D}_{\text{ofE}_1} \quad \mathcal{D}_{\text{ofE}'_1}}{t_1 \rightarrow t'_1 \quad \vdash t_1: \tau_1 \quad \vdash t'_1: \tau_1} \Rightarrow \quad \text{---} [\&_3] \\
\frac{\mathcal{D}_{\text{stepE}_1}}{\langle t_1, t_2 \rangle \rightarrow \langle t'_1, t_2 \rangle} [\text{step_pair1}] \quad \wedge \quad \frac{\mathcal{D}_{\text{ofE}_1} \quad \mathcal{D}_{\text{ofE}_2}}{\vdash t_1: \tau_1 \quad \vdash t_2: \tau_2} [\text{of_pair}] \quad \Rightarrow \quad \frac{\mathcal{D}_{\text{ofE}'_1} \quad \mathcal{D}_{\text{ofE}_2}}{\vdash \langle t'_1, t_2 \rangle: \tau_1 \times \tau_2} [\text{of_pair}]
\end{array}$$

$$\begin{array}{c}
\frac{\mathcal{D}_{\text{stepE}_2} \quad \mathcal{D}_{\text{ofE}_2} \quad \mathcal{D}_{\text{ofE}'_2}}{t_2 \rightarrow t'_2 \quad \vdash t_2: \tau_2 \quad \vdash t'_2: \tau_2} \Rightarrow \quad \text{---} [\&_4] \\
\frac{\mathcal{D}_{\text{valE}_1} \quad \mathcal{D}_{\text{stepE}_2}}{t_1 \text{ value} \quad t_2 \rightarrow t'_2} [\text{step_pair2}] \quad \wedge \quad \frac{\mathcal{D}_{\text{ofE}_1} \quad \mathcal{D}_{\text{ofE}_2}}{\vdash t_1: \tau_1 \quad \vdash t_2: \tau_2} [\text{of_pair}] \quad \Rightarrow \quad \frac{\mathcal{D}_{\text{ofE}_1} \quad \mathcal{D}_{\text{ofE}'_2}}{\vdash \langle t_1, t'_2 \rangle: \tau_1 \times \tau_2} [\text{of_pair}]
\end{array}$$

$$\begin{array}{c}
\frac{\mathcal{D}_{\text{val}_2} \quad \mathcal{D}_{\text{ofE}}}{\lambda x: \tau_2. t[x] \quad t_2 \rightarrow t[t_2]} [\text{step_beta}] \quad \wedge \quad \frac{\{x\} \vdash x: \tau_2 \rightarrow \vdash t[x]: \tau}{\vdash \lambda x: \tau_2. t[x]: \tau_2 \rightarrow \tau} [\text{of_lam}] \quad \frac{\mathcal{D}_{\text{ofE}_2}}{\vdash t_2: \tau_2} [\text{of_app}] \quad \Rightarrow \quad \frac{\mathcal{D}_{\text{ofE}} \quad t_2 \quad \mathcal{D}_{\text{ofE}_2}}{\vdash t[t_2]: \tau}
\end{array}$$

$$\frac{\mathcal{D}_{\text{val}_1} \quad \mathcal{D}_{\text{val}_2}}{t_1 \text{ \textbf{value}} \quad t_2 \text{ \textbf{value}}} \text{ [step_let]} \quad \wedge \quad \frac{\mathcal{D}_{\text{of}} \quad \frac{\mathcal{D}_{\text{of}_1} \quad \mathcal{D}_{\text{of}_2}}{\vdash t_1: \tau_1 \quad \vdash t_2: \tau_2} \text{ [of_pair]}}{\vdash \text{\textbf{let}} \langle x, y \rangle = \langle t_1, t_2 \rangle \text{ \textbf{in}} u[x][y]: \tau} \text{ [of_let]} \implies \frac{\mathcal{D}_{\text{of}} \quad t_1 \quad \mathcal{D}_{\text{of}_1} \quad t_2 \quad \mathcal{D}_{\text{of}_2}}{\vdash u[t_1][t_2]: \tau}$$

$$\begin{aligned}
\% \textbf{mode} \quad & +\mathcal{D}_{\text{step}} \quad \wedge \quad +\mathcal{D}_{\text{of}} \implies -\mathcal{D}_{\text{of}'} \\
\% \textbf{worlds} \quad & () \quad \mathcal{D}_{\text{step}} \quad \wedge \quad \mathcal{D}_{\text{of}} \implies \mathcal{D}_{\text{of}'} \\
\% \textbf{total} \quad & \mathcal{D}_{\text{step}} \quad \mathcal{D}_{\text{step}} \quad \wedge \quad \mathcal{D}_{\text{of}} \implies \mathcal{D}_{\text{of}'}
\end{aligned}$$