

1 System T

1.1 Syntax

$term$: **type**.
 $clos$: **type**.
 env : **type**.
 $stack$: **type**.
 $state$: **type**.
 int : **type**.
 $bool$: **type**.
 $frame$: **type**.

Term t

x : $term$.
 m : $term$.
 $t_1 t_2$: $term$.
 $\lambda x.t$: $term$.
 $\mathbf{succ}(t)$: $term$.
 $\mathbf{pred}(t)$: $term$.
 $\mathbf{rec}(t_1, t_2, t_3)$: $term$.
 $\mathbf{let } x = t_1 \mathbf{ in } t_2$: $term$.

Integer m

\mathbf{Z} : int .
 $\mathbf{S}(m)$: int .

Boolean b

\mathbf{T} : $bool$.
 \mathbf{F} : $bool$.

Closure c

$(t; \mathcal{E})$: $clos$.

Environment \mathcal{E}

\square : env .
 $(\mathcal{E}, x \leftarrow c)$: env .

Frame f

$(\square c)$: $frame$.
 $(c \square)$: $frame$.
 $\mathbf{rec}(\square, c_2, c_3)$: $frame$.
 $\mathbf{rec}(c_1, \square, c_3)$: $frame$.
 $\mathbf{rec}(c_1, c_2, \square)$: $frame$.
 $\mathbf{succ}(\square)$: $frame$.

Stack \mathcal{S}

\square : $stack$.

$f:\mathcal{S} \quad :$ *stack*.

State σ

$\langle t, \mathcal{E}, \mathcal{S} \rangle \quad :$ *state*.

1.2 Judgments

$\mathcal{E}(x)=c \quad :$ **type**.

$\sigma 1 \rightarrow \sigma 2 \quad :$ **type**.

$\sigma 1 \rightsquigarrow c_2 \quad :$ **type**.

$t \text{ value} \quad :$ **type**.

1.3 Fetch

$$\overline{(\mathcal{E}, x \leftarrow c)(x)=c} \quad [\text{Fetch1}]$$

$$\frac{x \neq x' \quad \mathcal{E}(x)=c}{(\mathcal{E}, x' \leftarrow c')(x)=c} \quad [\text{Fetch2}]$$

%mode $+ \mathcal{E}(+x) = -c$

%worlds $() \quad \mathcal{E}(x)=c$

%terminates $\mathcal{E} \quad \mathcal{E}(x)=c$

Remark. Twelf cannot check the following property:

%unique $+ \mathcal{E}(+x) = -1c$

1.4 Value

$$\overline{m \text{ value}} \quad [\text{V_Cst}]$$

$$\overline{x \text{ value}} \quad [\text{V_Var}]$$

$$\overline{\lambda x.t \text{ value}} \quad [\text{V_Abs}]$$

1.5 Evaluation

$$\frac{\mathcal{E}(x)=(t; \mathcal{E}')}{\langle x, \mathcal{E}, \mathcal{S} \rangle \rightarrow \langle t, \mathcal{E}', \mathcal{S} \rangle} \quad [\text{E_Var}]$$

$$\overline{\langle t t', \mathcal{E}, \mathcal{S} \rangle \rightarrow \langle t', \mathcal{E}, ((t; \mathcal{E}) \square) : \mathcal{S} \rangle} \quad [\text{E_App1}]$$

$$\frac{w \text{ value}}{\langle w, \mathcal{E}, (((t; \mathcal{E}') \square)) : \mathcal{S} \rangle \rightarrow \langle t, \mathcal{E}', (\square (w; \mathcal{E})) : \mathcal{S} \rangle} \quad [\text{E_App2}]$$

$$\overline{\langle \lambda x.t, \mathcal{E}, (\square c) : \mathcal{S} \rangle \rightarrow \langle t, (\mathcal{E}, x \leftarrow c), \mathcal{S} \rangle} \quad [\text{E_Abs}]$$

$$\overline{\langle \text{let } x = t_1 \text{ in } t_2, \mathcal{E}, \mathcal{S} \rangle \rightarrow \langle t_1, \mathcal{E}, ((\lambda x.t_2; \mathcal{E}) \square) : \mathcal{S} \rangle} \quad [\text{E_Let}]$$

$$\overline{\langle \text{succ}(t), \mathcal{E}, \mathcal{S} \rangle \rightarrow \langle t, \mathcal{E}, \text{succ}(\square) : \mathcal{S} \rangle} \quad [\text{E_succ1}]$$

$$\overline{\langle m, \mathcal{E}, \text{succ}(\square) : \mathcal{S} \rangle \rightarrow \langle \mathbf{S}(m), \mathcal{E}, \mathcal{S} \rangle} \quad [\text{E_succ2}]$$

$$\overline{\langle \text{rec}(t_1, t_2, t_3), \mathcal{E}, \mathcal{S} \rangle \rightarrow \langle t_3, \mathcal{E}, \text{rec}((t_1; \mathcal{E}), (t_2; \mathcal{E}), \square) : \mathcal{S} \rangle} \quad [\text{E_rec3}]$$

$$\begin{array}{c}
\frac{w \text{ value}}{\langle w, \mathcal{E}, \text{rec}((t_1; \mathcal{E}_1), (t_2; \mathcal{E}_2), []) : \mathcal{S} \rangle \rightarrow \langle t_2, \mathcal{E}_2, \text{rec}((t_1; \mathcal{E}), [], (w; \mathcal{E})) : \mathcal{S} \rangle} [\text{E_rec4}] \\
\\
\frac{w \text{ value} \quad w_3 \text{ value}}{\langle w, \mathcal{E}, \text{rec}((t_1; \mathcal{E}_1), [], (w_3; \mathcal{E}_3)) : \mathcal{S} \rangle \rightarrow \langle t_1, \mathcal{E}_1, \text{rec}([], (w; \mathcal{E}), (w_3; \mathcal{E}_3)) : \mathcal{S} \rangle} [\text{E_rec5}] \\
\\
\frac{w_2 \text{ value} \quad w_3 \text{ value}}{\langle \mathbf{Z}, \mathcal{E}, \text{rec}([], (w_2; \mathcal{E}_2), (w_3; \mathcal{E}_3)) : \mathcal{S} \rangle \rightarrow \langle w_2, \mathcal{E}_2, \mathcal{S} \rangle} [\text{E_rec1}] \\
\\
\frac{w_2 \text{ value} \quad w_3 \text{ value}}{\langle \mathbf{S}(m), \mathcal{E}, \text{rec}([], (w_2; \mathcal{E}_2), (w_3; \mathcal{E}_3)) : \mathcal{S} \rangle \rightarrow \langle m, \mathcal{E}, \text{rec}([], (w_2; \mathcal{E}_2), (w_3; \mathcal{E}_3)) : ((w_3 m); \mathcal{E}_3) [] : \mathcal{S} \rangle} [\text{E_rec2}]
\end{array}$$

```

%mode   +t value
%worlds () t value
%mode   +σ1 → -σ2
%worlds () σ1 → σ2

```

Remark. Twelf cannot check the following property:

```
%unique +σ1 → -1σ2
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1.6 Full evaluation

$$\begin{array}{c}
\frac{m \text{ value}}{\langle m, \mathcal{E}, [] \rangle \rightsquigarrow (m; \mathcal{E})} [\text{Eval_1}] \\
\\
\frac{\sigma_1 \rightarrow \sigma_2 \quad \sigma_2 \rightsquigarrow c}{\sigma_1 \rightsquigarrow c} [\text{Eval_2}] \\
\\
\begin{array}{l}
\text{\%solve} \quad \langle \text{succ}(\mathbf{Z}), [], [] \rangle \rightsquigarrow \sqcup \\
\text{\%solve} \quad \langle \text{let } g = \lambda x. \text{succ}(x) \text{ in } (g \mathbf{Z}), [], [] \rangle \rightsquigarrow \sqcup
\end{array}
\end{array}$$