

# 1 Equality

## 1.1 Syntax

### 1.1.1 Types

**%datatype** *type*  
**%name** *type*  $\tau$

$\tau ::=$   
 $\quad | \text{unit}$   
 $\quad | \tau_1 \rightarrow \tau_2$

### 1.1.2 Terms

**%datatype** *term*  
**%name** *term*  $t$

$t ::=$   
 $\quad | \langle \rangle$   
 $\quad | t_1 t_2$   
 $\quad | \lambda x: \tau. t$

**%binding**  $1 \mapsto 3$  **in**  $\lambda \square: \square. \square$

## 1.2 Typing judgment

**%judgment**  $\vdash t: \tau$

$$\frac{}{\vdash \langle \rangle: \mathbf{unit}} [\text{of\_empty}]$$

$$\frac{\{x\} \vdash x: \tau \rightarrow \vdash t[x]: \tau'}{\vdash \lambda x: \tau. t[x]: \tau \rightarrow \tau'} [\text{of\_lam}]$$

$$\frac{\vdash t_1: \tau \rightarrow \tau' \quad \vdash t_2: \tau}{\vdash t_1 t_2: \tau'} [\text{of\_app}]$$

## 1.3 Values

**%judgment**  $t$  **value**

$$\langle \rangle \text{ value} [\text{value\_empty}]$$

$$\lambda x: \tau. t[x] \text{ value} [\text{value\_lam}]$$

## 1.4 Algorithmic definition of equality

**%judgment**  $t_1 = t_2$

$$\begin{array}{c}
\overline{\langle \rangle = \langle \rangle}^{\text{[equ\_u]}} \\
\frac{\{x\} \ x = x \rightarrow t[x] = t'[x]}{\lambda x: \tau.t[x] = \lambda x: \tau.t'[x]}^{\text{[equ\_l]}} \\
\frac{t_1 = e_1 \quad t_2 = e_2}{(t_1 t_2) = (e_1 e_2)}^{\text{[equ\_a]}}
\end{array}$$

### 1.5 Reflexivity is admissible (first attempt)

**%lemma**  $\forall t: \text{term} \quad \vdash_{\text{ref}} \exists \mathcal{D}: t = t$   
**%mode**  $+t \quad \vdash_{\text{ref}} \quad -\mathcal{D}$

$$\begin{array}{c}
\overline{\langle \rangle: \text{term} \quad \vdash_{\text{ref}} \quad \overline{\langle \rangle = \langle \rangle}^{\text{[equ\_u]}}}^{\text{[r\_a]}} \\
\\
\frac{\{x\} \ \{\mathcal{U}: x = x\} \left( t[x]: \text{term} \quad \vdash_{\text{ref}} \quad \frac{\mathcal{D} \ x \ \mathcal{U}}{t[x] = t[x]} \right)}{\mathcal{D}}^{\text{[r\_l]}} \\
\lambda x: \tau.t[x]: \text{term} \quad \vdash_{\text{ref}} \quad \frac{\{x\} \ x = x \rightarrow t[x] = t[x]}{\lambda x: \tau.t[x] = \lambda x: \tau.t[x]}^{\text{[equ\_l]}} \\
\\
\frac{\left( t_1: \text{term} \quad \vdash_{\text{ref}} \quad \frac{\mathcal{D}_1}{t_1 = t_1} \right) \left( t_2: \text{term} \quad \vdash_{\text{ref}} \quad \frac{\mathcal{D}_2}{t_2 = t_2} \right)}{\frac{\mathcal{D}_1 \quad \mathcal{D}_2}{(t_1 t_2): \text{term} \quad \vdash_{\text{ref}} \quad \frac{t_1 = t_1 \quad t_2 = t_2}{(t_1 t_2) = (t_1 t_2)}}^{\text{[equ\_a]}}}^{\text{[r\_a]}}
\end{array}$$

**%block**  $\mathcal{W}_{\text{ref}} : \text{block } \{x: \text{term}\} \{\mathcal{U}: x = x\}$   
**%worlds**  $(\mathcal{W}_{\text{ref}}) \ t \quad \vdash_{\text{ref}} \ \mathcal{D}$   
**%terminates**  $(t) \ t \quad \vdash_{\text{ref}} \ \mathcal{D}$

**Remark.** Checking this totality assertion yields “Coverage error”:

**%total**  $(t) \ t \quad \vdash_{\text{ref}} \ \mathcal{D}$

### 1.6 Reflexivity is admissible

**%lemma**  $\forall t: \text{term} \quad \vdash_{\text{ref}} \exists \mathcal{D}: t = t$   
**%mode**  $+t \quad \vdash_{\text{ref}} \quad -\mathcal{D}$

$$\begin{array}{c}
\overline{\langle \rangle: \text{term} \quad \vdash_{\text{ref}} \quad \overline{\langle \rangle = \langle \rangle}^{\text{[equ\_u]}}}^{\text{[r\_u]}} \\
\\
\frac{\{x\} \ \{\mathcal{U}: x = x\} \left( x: \text{term} \quad \vdash_{\text{ref}} \quad \frac{\mathcal{U}}{x = x} \right) \rightarrow \left( t[x]: \text{term} \quad \vdash_{\text{ref}} \quad \frac{\mathcal{D} \ x \ \mathcal{U}}{t[x] = t[x]} \right)}{\mathcal{D}}^{\text{[r\_l]}} \\
\lambda x: \tau.t[x]: \text{term} \quad \vdash_{\text{ref}} \quad \frac{\{x\} \ x = x \rightarrow t[x] = t[x]}{\lambda x: \tau.t[x] = \lambda x: \tau.t[x]}^{\text{[equ\_l]}}
\end{array}$$

$$\frac{\left( \frac{t_1: term \quad \vdash_{\text{ref}} \quad \mathcal{D}_1}{t_1 = t_1} \right) \quad \left( \frac{t_2: term \quad \vdash_{\text{ref}} \quad \mathcal{D}_2}{t_2 = t_2} \right)}{\frac{\mathcal{D}_1 \quad \mathcal{D}_2}{(t_1 t_2): term \quad \vdash_{\text{ref}} \quad \frac{t_1 = t_1 \quad t_2 = t_2}{(t_1 t_2) = (t_1 t_2)} [\text{equ\_a}]} [\text{r\_a}]}$$

**%block**  $\mathcal{W}_{\text{ref}} : \mathbf{block} \{x: term\} \{ \mathcal{U}: x = x \} \{ \sqcup : (x: term \quad \vdash_{\text{ref}} \quad \mathcal{U}: x = x) \}$   
**%worlds**  $(\mathcal{W}_{\text{ref}}) \quad t \quad \vdash_{\text{ref}} \quad \mathcal{D}$   
**%terminates**  $(t) \quad t \quad \vdash_{\text{ref}} \quad \mathcal{D}$   
**%total**  $(t) \quad t \quad \vdash_{\text{ref}} \quad \mathcal{D}$

## 1.7 Transitivity is admissible

**%lemma**  $t = r \quad \wedge \quad r = s \quad \vdash_{\text{tr}} \quad t = s$   
**%mode**  $+\mathcal{E}_1 \quad \wedge \quad +\mathcal{E}_2 \quad \vdash_{\text{tr}} \quad -\mathcal{E}_3$

$$\frac{\frac{\frac{}{\overline{\langle \rangle = \langle \rangle} [\text{equ\_u}]} \quad \wedge \quad \frac{}{\overline{\langle \rangle = \langle \rangle} [\text{equ\_u}]} \quad \vdash_{\text{tr}} \quad \frac{}{\overline{\langle \rangle = \langle \rangle} [\text{equ\_u}]} [\text{tr\_u}]}{\frac{\mathcal{E}_1}{\lambda x: \tau. t[x] = \lambda x: \tau. r[x]} [\text{equ\_l}]} \quad \wedge \quad \frac{\mathcal{E}_2}{\lambda x: \tau. r[x] = \lambda x: \tau. s[x]} [\text{equ\_l}] \quad \vdash_{\text{tr}} \quad \frac{\mathcal{E}_3}{\lambda x: \tau. t[x] = \lambda x: \tau. s[x]} [\text{equ\_l}]} [\text{tr\_l}]}$$

$$\frac{\left( \frac{\mathcal{E}_1}{t_1 = r_1} \quad \wedge \quad \frac{\mathcal{F}_1}{r_1 = s_1} \quad \vdash_{\text{tr}} \quad \frac{\mathcal{G}_1}{t_1 = s_1} \right) \quad \left( \frac{\mathcal{E}_2}{t_2 = r_2} \quad \wedge \quad \frac{\mathcal{F}_2}{r_2 = s_2} \quad \vdash_{\text{tr}} \quad \frac{\mathcal{G}_2}{t_2 = s_2} \right)}{\frac{\mathcal{E}_1 \quad \mathcal{E}_2}{\frac{t_1 = r_1 \quad t_2 = r_2}{(t_1 t_2) = (r_1 r_2)} [\text{equ\_a}]} \quad \wedge \quad \frac{\mathcal{F}_1 \quad \mathcal{F}_2}{\frac{r_1 = s_1 \quad r_2 = s_2}{(r_1 r_2) = (s_1 s_2)} [\text{equ\_a}]} \quad \vdash_{\text{tr}} \quad \frac{\mathcal{G}_1 \quad \mathcal{G}_2}{\frac{t_1 = s_1 \quad t_2 = s_2}{(t_1 t_2) = (s_1 s_2)} [\text{equ\_a}]} [\text{tr\_a}]}$$

**%block**  $\mathcal{W}_{\text{tr}} : \mathbf{block} \{x: term\} \{ \mathcal{U}: x = x \} \{ \sqcup : (\mathcal{U}: x = x \quad \wedge \quad \mathcal{U}: x = x \quad \vdash_{\text{tr}} \quad \mathcal{U}: x = x) \}$   
**%worlds**  $(\mathcal{W}_{\text{tr}}) \quad \mathcal{E}_1 \quad \wedge \quad \mathcal{E}_2 \quad \vdash_{\text{tr}} \quad \mathcal{E}_3$   
**%terminates**  $(\mathcal{E}_1) \quad \mathcal{E}_1 \quad \wedge \quad \mathcal{E}_2 \quad \vdash_{\text{tr}} \quad \mathcal{E}_3$   
**%total**  $(\mathcal{E}_1) \quad \mathcal{E}_1 \quad \wedge \quad \mathcal{E}_2 \quad \vdash_{\text{tr}} \quad \mathcal{E}_3$

## 1.8 Symmetry is admissible

**%lemma**  $t = r \quad \vdash_{\text{sym}} \quad r = t$   
**%mode**  $+\mathcal{E}_1 \quad \vdash_{\text{sym}} \quad -\mathcal{E}_2$

$$\frac{\frac{\frac{}{\overline{\langle \rangle = \langle \rangle} [\text{equ\_u}]} \quad \vdash_{\text{sym}} \quad \frac{}{\overline{\langle \rangle = \langle \rangle} [\text{equ\_u}]} [\text{tr\_u}]}{\frac{\mathcal{E}_1}{\{x\} \quad x = x \rightarrow t[x] = r[x]} [\text{equ\_l}]} \quad \vdash_{\text{sym}} \quad \frac{\mathcal{E}'_1}{\{x\} \quad x = x \rightarrow r[x] = t[x]} [\text{equ\_l}]} [\text{tr\_l}]}$$

$$\frac{\left( \frac{\mathcal{E}_1}{t_1 = r_1} \vdash_{\text{sym}} \frac{\mathcal{E}'_1}{r_1 = t_1} \right) \left( \frac{\mathcal{E}_2}{t_2 = r_2} \vdash_{\text{sym}} \frac{\mathcal{E}'_2}{r_2 = t_2} \right)}{\frac{\mathcal{E}_1 \quad \mathcal{E}_2}{t_1 = r_1 \quad t_2 = r_2} \text{[equ\_a]} \vdash_{\text{sym}} \frac{\mathcal{E}'_1 \quad \mathcal{E}'_2}{r_1 = t_1 \quad r_2 = t_2} \text{[equ\_a]}} \text{[tr\_a]}$$

**%block**  $\mathcal{W}_{\text{sym}} : \mathbf{block} \{x: \text{term}\} \{\mathcal{U}: x = x\} \{\sqcup: (\mathcal{U}: x = x \vdash_{\text{sym}} \mathcal{U}: x = x)\}$   
**%worlds**  $(\mathcal{W}_{\text{sym}}) \mathcal{E}_1 \vdash_{\text{sym}} \mathcal{E}_2$   
**%terminates**  $(\mathcal{E}_1) \mathcal{E}_1 \vdash_{\text{sym}} \mathcal{E}_2$   
**%total**  $(\mathcal{E}_1) \mathcal{E}_1 \vdash_{\text{sym}} \mathcal{E}_2$