

1 First simulation (soundness)

1.1 Syntax

1.1.1 Index, indices and tables

%datatype *index*

%name *index* $n \ \alpha$

$n ::= 0$
 $\quad \mid n+1$

%datatype *vector*

%name *vector* \mathcal{I}

$\mathcal{I} ::= []$
 $\quad \mid n :: \mathcal{I}$

%datatype *table*

%name *table* \mathcal{I}_α

$\mathcal{I}_\alpha ::= []$
 $\quad \mid \mathcal{I} :: \mathcal{I}_\alpha$

1.1.2 Term

%datatype *term*

%name *term* t

$t ::= n$
 $\quad \mid t_1 t_2$
 $\quad \mid \lambda t$
 $\quad \mid \text{catch } t$
 $\quad \mid \text{throw } \alpha \ t$

Remark. Syntax of safe λ_{ct} -terms:

$t ::= n$
 $\quad \mid t_1 t_2$
 $\quad \mid \lambda t$
 $\quad \mid \text{get-context } t$
 $\quad \mid \text{set-context } \alpha \ t$

1.2 Subtraction

%judgment $n_1 \dot{-} n_2 = n_3$

$n_1 \dot{-} 0 = n_1$ [minus]

$(n_1 + 1) \dot{-} (n_2 + 1) = n_3$ [minus] *when* $n_1 \dot{-} n_2 = n_3$

%mode $+n_1 \dot{-} +n_2 = -n_3$

%worlds $() \ n_1 \dot{-} n_2 = n_3$

%terminates $(n_1) \ n_1 \dot{-} n_2 = n_3$

%unique $+n_1 \dot{-} +n_2 = -1n_3$

%lemma $\forall n: \text{index} \cdot n \dot{-} n = 0$ [minus id]

Proof.

$$\frac{}{0: \text{index} \cdot 0 \dot{-} 0 = 0} \text{ [minus-0]} \quad \text{[minus-id]} \quad [k1]$$

$$\frac{n: \text{index} \cdot n \dot{-} n = 0 \quad \text{[minus id]}}{n+1: \text{index} \cdot (n+1) \dot{-} (n+1) = 0} \text{ [minus-id]} \quad [k2]$$

%mode $+n \cdot -\mathcal{D}$ [minus id]

%worlds $() \ n \cdot \mathcal{D}$ [minus id]

%total $(\alpha) \ n \cdot \mathcal{D}$ [minus id]

%lemma $n_1 \dot{-} (n_2 + 1) = n_3 \Rightarrow n_1 \dot{-} n_2 = (n_3 + 1)$ [minus-acc]

Proof.

$$\frac{}{(n_1 + 1) \dot{-} (0 + 1) = n_1} \text{ [minus-acc]} \Rightarrow (n_1 + 1) \dot{-} 0 = (n_1 + 1) \text{ [minus-id]} \quad \text{[minus-acc]} \quad [k1]$$

$$\frac{\frac{\mathcal{D}_1}{n_1 \dot{-} (n_2 + 1) = n_3} \Rightarrow \frac{\mathcal{D}_2}{n_1 \dot{-} n_2 = (n_3 + 1)} \text{ [minus-acc]} \quad \frac{\mathcal{D}_1}{(n_1 + 1) \dot{-} ((n_2 + 1) + 1) = n_3} \Rightarrow \frac{\mathcal{D}_2}{(n_1 + 1) \dot{-} (n_2 + 1) = (n_3 + 1)} \text{ [minus-id]} \quad \text{[minus-acc]} \quad [k2]$$

%mode $+\mathcal{D}_1 \Rightarrow -\mathcal{D}_2$ [minus-acc]

%worlds $() \ \mathcal{D}_1 \Rightarrow \mathcal{D}_2$ [minus-acc]

%total $\mathcal{D}_1 \ \mathcal{D}_1 \Rightarrow \mathcal{D}_2$ [minus-acc]

%lemma $n_1 \dot{-} n_2 = n_3 \Rightarrow n_1 \dot{-} n_3 = n_2$ [minus-swap]

Proof.

$$\frac{n_1: \text{index} \cdot n_1 \dot{-} n_1 = 0 \quad \text{[minus id]}}{n_1 \dot{-} 0 = n_1} \text{ [minus]} \Rightarrow n_1 \dot{-} n_1 = 0 \quad \text{[minus-swap]} \quad [k1]$$

$$\frac{\frac{\mathcal{D}_1}{n_1 \dot{-} n_2 = n_3} \Rightarrow \frac{\mathcal{D}_2}{n_1 \dot{-} n_3 = n_2} \text{ [minus-swap]} \quad \frac{\mathcal{D}_2}{n_2 \dot{-} n_3 = n_2} \Rightarrow \frac{\mathcal{D}_3}{(n_1 + 1) \dot{-} (n_2 + 1) = n_3} \Rightarrow \frac{\mathcal{D}_3}{(n_1 + 1) \dot{-} n_3 = (n_2 + 1)} \text{ [minus-acc]} \quad \frac{\mathcal{D}_1}{(n_1 + 1) \dot{-} n_2 = n_3} \Rightarrow \frac{\mathcal{D}_3}{(n_1 + 1) \dot{-} (n_2 + 1) = n_3} \text{ [minus-id]} \quad \text{[minus-swap]} \quad [k2]$$

%mode $+\mathcal{D}_1 \Rightarrow -\mathcal{D}_2$ [minus-swap]

%worlds $() \ \mathcal{D}_1 \Rightarrow \mathcal{D}_2$ [minus-swap]

%total $\mathcal{D}_1 \ \mathcal{D}_1 \Rightarrow \mathcal{D}_2$ [minus-swap]

1.2.1 Fetch (indices)

%judgment $\mathcal{I}(n_1) = n_2$

$(n :: \mathcal{I})(0) = n$ [fetch]

$(n :: \mathcal{I})(n_1 + 1) = n_2$ [fetch] *when* $\mathcal{I}(n_1) = n_2$

%mode $+ \mathcal{I}(+n_1) = -n_2$

%worlds $() \ \mathcal{I}(n_1) = n_2$

%terminates $n_1 \ \mathcal{I}(n_1) = n_2$

%unique $+ \mathcal{I}(+n_1) = -1n_2$

1.2.2 Fetch (table)

%judgment $\mathcal{I}_\alpha(n) = \mathcal{I}$

$(\mathcal{I} :: \mathcal{I}_\alpha)(0) = \mathcal{I}$ [fetch₁]

$(\mathcal{I}' :: \mathcal{I}_\alpha)(\alpha + 1) = \mathcal{I}'$ [fetch₂] *when* $\mathcal{I}_\alpha(\alpha) = \mathcal{I}$

%mode $+ \mathcal{I}_\alpha(+\alpha) = -\mathcal{I}$

%worlds $() \ \mathcal{I}_\alpha(\alpha) = \mathcal{I}$

%terminates $\alpha \ \mathcal{I}_\alpha(\alpha) = \mathcal{I}$

%unique $+ \mathcal{I}_\alpha(+\alpha) = -1\mathcal{I}$

1.2.3 Compute

%judgment $n_1 \dot{-} \mathcal{I}(n_2) = n_3$

$n \dot{-} \mathcal{I}(l) = g$ [compute] *when* $\mathcal{I}(l) = k, \ n \dot{-} k = g$

%mode $+n_1 \dot{-} +(\mathcal{I}(+n_2)) = -n_3$

%worlds $() \ n_1 \dot{-} \mathcal{I}(n_2) = n_3$

%terminates $\{\} \ n_1 \dot{-} \mathcal{I}(n_2) = n_3$

%unique $+n_1 \dot{-} +(\mathcal{I}(+n_2)) = -1n_3$

2 Safe λ_{ct} -terms

2.1 Safety

%judgment $n \in \mathcal{I}$

$n \in (n :: \mathcal{I})$ [members]

$n \in (n' :: \mathcal{I})$ [members] *when* $n \in \mathcal{I}$

%mode $+n \in +\mathcal{I}$

%worlds $() \ n \in \mathcal{I}$

%terminates $\mathcal{I} \ n \in \mathcal{I}$

%lemma $\mathcal{I}(n) = k \Rightarrow k \in \mathcal{I}$ [target]

Proof.

$$\frac{}{(n :: \mathcal{I})(0) = n} \text{ [fetch]} \Rightarrow n \in (n :: \mathcal{I}) \text{ [members]} \quad \text{[target]} \quad [k1]$$

$$\frac{\frac{\mathcal{D}_1}{\mathcal{I}(n) = k} \Rightarrow \frac{\mathcal{D}_2}{k \in \mathcal{I}} \text{ [target]} \quad \frac{\mathcal{D}_1}{\mathcal{I}(n) = k} \Rightarrow \frac{\mathcal{D}_2}{(k' :: \mathcal{I})(n+1) = k'} \text{ [fetch]} \Rightarrow \frac{\mathcal{D}_2}{k \in (k' :: \mathcal{I})} \text{ [members]} \quad \text{[target]} \quad [k2]$$

%mode $+\mathcal{D}_1 \Rightarrow -\mathcal{D}_2$ [target]

%worlds $() \ \mathcal{D}_1 \Rightarrow \mathcal{D}_2$ [target]

%total $(\mathcal{D}_1) \ \mathcal{D}_1 \Rightarrow \mathcal{D}_2$ [target]

%judgment $\text{Safe}_n^{\mathcal{I}, \mathcal{I}_\alpha, \mathcal{I}_\beta}(t)$

$\text{Safe}_n^{\mathcal{I}, \mathcal{I}_\alpha}(g)$ [safe] *when* $n \dot{-} g = k, \ k \in \mathcal{I}$

$\text{Safe}_n^{\mathcal{I}, \mathcal{I}_\alpha}(t \ u)$ [safe] *when* $\text{Safe}_n^{\mathcal{I}, \mathcal{I}_\alpha}(t), \ \text{Safe}_n^{\mathcal{I}, \mathcal{I}_\alpha}(u)$

$\text{Safe}_n^{\mathcal{I}, \mathcal{I}_\alpha}(\lambda t)$ [safe] *when* $\text{Safe}_{n+1}^{\mathcal{I}, \mathcal{I}_\alpha}(\mathcal{I}_\alpha(t)) = t'$

$\text{Safe}_n^{\mathcal{I}, \mathcal{I}_\alpha}(\text{catch } t)$ [safe] *when* $\text{Safe}_n^{\mathcal{I}, \mathcal{I}_\alpha}(\mathcal{I}_\alpha(t)) = t'$

$\text{Safe}_n^{\mathcal{I}, \mathcal{I}_\alpha}(\text{throw } \alpha \ t)$ [safe] *when* $\mathcal{I}_\alpha(\alpha) = \mathcal{I}', \ \text{Safe}_n^{\mathcal{I}', \mathcal{I}_\alpha}(t) = t'$

%mode $\text{Safe}_n^{\mathcal{I}, \mathcal{I}_\alpha, \mathcal{I}_\beta}(+t) = -t'$

%worlds $() \ \text{Safe}_n^{\mathcal{I}, \mathcal{I}_\alpha, \mathcal{I}_\beta}(t) = t'$

%terminates $t \ \text{Safe}_n^{\mathcal{I}, \mathcal{I}_\alpha, \mathcal{I}_\beta}(t) = t'$

%unique $\text{Safe}_n^{\mathcal{I}, \mathcal{I}_\alpha, \mathcal{I}_\beta}(+t) = -1t'$

2.2 From local indices to global indices

%judgment $\dot{-}_n^{\mathcal{I}, \mathcal{I}_\alpha}(t_1) = t_2$

$\dot{-}_n^{\mathcal{I}, \mathcal{I}_\alpha}(l) = g$ [id] *when* $n \dot{-} \mathcal{I}(l) = g$

$\dot{-}_n^{\mathcal{I}, \mathcal{I}_\alpha}(t \ u) = t' \ u'$ [id] *when* $\dot{-}_n^{\mathcal{I}, \mathcal{I}_\alpha}(t) = t', \ \dot{-}_n^{\mathcal{I}, \mathcal{I}_\alpha}(u) = u'$

$\dot{-}_n^{\mathcal{I}, \mathcal{I}_\alpha}(\lambda t) = \lambda t'$ [id] *when* $\dot{-}_{n+1}^{\mathcal{I}, \mathcal{I}_\alpha}(\mathcal{I}_\alpha(t)) = t'$

$\dot{-}_n^{\mathcal{I}, \mathcal{I}_\alpha}(\text{get-context } t) = \text{catch } t'$ [id] *when* $\dot{-}_n^{\mathcal{I}, \mathcal{I}_\alpha}(\mathcal{I}_\alpha(t)) = t'$

$\dot{-}_n^{\mathcal{I}, \mathcal{I}_\alpha}(\text{set-context } \alpha \ t) = \text{throw } \alpha \ t'$ [id] *when* $\mathcal{I}_\alpha(\alpha) = \mathcal{I}', \ \dot{-}_n^{\mathcal{I}', \mathcal{I}_\alpha}(t) = t'$

%mode $\dot{-}_n^{\mathcal{I}, \mathcal{I}_\alpha, \mathcal{I}_\beta}(+t) = -t'$

%worlds $() \ \dot{-}_n^{\mathcal{I}, \mathcal{I}_\alpha, \mathcal{I}_\beta}(t) = t'$

%terminates $t \ \dot{-}_n^{\mathcal{I}, \mathcal{I}_\alpha, \mathcal{I}_\beta}(t) = t'$

%unique $\dot{-}_n^{\mathcal{I}, \mathcal{I}_\alpha, \mathcal{I}_\beta}(+t) = -1t'$

