# Partial Lagrangian relaxation for General Quadratic Programming 

Alain Faye, Frédéric Roupin<br>CEDRIC, CNAM-IIE

SCRO/JOPT May 9, 2006

Main results (to appear in 4OR)

1. A complete characterization of constant quadratic functions over an affine variety $\Omega=\left\{x \in \Re^{n}: A x=b\right\}$.
2. How to convexify the objective function of a general quadratic programming problem (Pb) by using the linear constraints.
3. Formulation as a semidefinite program of the partial Lagrangian relaxation of ( Pb ) where the linear constraints are not relaxed.
4. Comparison of two semidefinite relaxations made from two sets of null quadratic functions over an affine variety.

## General Quadratic Programming

$$
\text { (Pb) } \min _{x} f(x)=x^{t} Q x+c^{t} x \text { s.t. } \begin{cases}x^{t} B_{i} x+d_{i}^{t} x=e_{i} & i \in I= \\ x^{t} B_{i} x+d_{i}^{t} x \leq e_{i} & i \in I \leq \\ A x=b & \end{cases}
$$

$\diamond$ Boolean quadratic problems can be formulated as (Pb) by considering $x_{i}^{2}=x_{i}$ for all $i$ in $\{1, \ldots, n\}$
$\diamond$ Standard semidefinite relaxations of $0-1$ quadratic programs are nothing but particular instances of Lagrangian duality [Poljak et al,1995] [Lemaréchal,2003]
$\diamond$ For the Boolean case, the supremum of the corresponding augmented Lagrangian function is equal to the value of the partial Lagrangian dual (the linear constraints are not relaxed) [Lemaréchal, Oustry, 2001]

Characterization of constant quadratic functions over $\Omega=\left\{x \in \Re^{n}: A x=b\right\}$

Let $F(x)=x^{t} H x+g^{t} x$ be a quadratic function that takes a constant value over $\Omega$. (underlying idea: add redundant constraints)

For any $x \in \Omega, u \in \operatorname{Ker}(A)=\left\{u \in \Re^{n}: A u=0\right\}$, and $\lambda \in \Re$
$F(x+\lambda u)=F(x)=F(x)+\lambda^{2} u^{t} H u+2 \lambda u^{t} H x+\lambda g^{t} u$
i.e. $F(x+\lambda u)$ does not depend on $\lambda$.

Necessary and sufficient conditions on $F(x)$ :

$$
\begin{aligned}
u^{t} H u & =0 \forall u \in \operatorname{Ker}(A) \\
2 u^{t} H x+g^{t} u & =0 \forall u \in \operatorname{Ker}(A) \forall x \in \Omega
\end{aligned} \text { [B] }
$$

## Characterization of constant quadratic functions

over $\Omega=\left\{x \in \Re^{n}: A x=b\right\}$

Lemma. $q(u)=u^{t} H u$ is a null quadratic form over $\operatorname{Ker}(A)$ if and only if $q(u)=u^{t}\left(A^{t} W^{t}+W A\right) u$, where $W$ is any $n \times p$-matrix.

Sketch proof. choose a "good" basis for the quadratic form, write and simplify $P^{t} H P$, then obtain $H=\left(P^{t}\right)^{-1} P^{t} H P P^{-1}$. $P=\left[\begin{array}{ll}A^{t} & B\end{array}\right]$, the $n-p$ columns of $B$ are a basis of $\operatorname{Ker}(A)$.

Theorem. $F(x)=x^{t} H x+g^{t} x$ is a constant quadratic function over $\{x: A x=b\}$ if and only if $F(x)=x^{t}\left(A^{t} W^{t}+W A\right) x+$ $\left(A^{t} \alpha-2 W b\right)^{t} x$, where $W$ is a $n \times p$-matrix and $\alpha$ is a $p$-vector.

Any constant quadratic function over $\Omega=\left\{x \in \Re^{n}: A x=b\right\}$ can be obtained by setting particular values to $W$ and $\alpha$ in

$$
F(x)=x^{t}\left(A^{t} W^{t}+W A\right) x+\left(A^{t} \alpha-2 W b\right)^{t} x
$$

$\diamond$ Product member by variable
$W=\frac{1}{2} E^{i j}$ and $\alpha=0$ where $E_{i j}^{i j}=1$ otherwise $E_{i j}^{k l}=0$.
$F(x)=x_{i} a_{j}^{t} x-b_{j} x_{i}=0 \forall x \in \Omega$
$\diamond$ Product member by member
$W=\frac{1}{2} A^{t} V$ with $V=\frac{1}{2}\left(E^{i j}+E^{j i}\right) F(x)=x^{t} a_{i} a_{j}^{t} x+x^{t} A^{t}(\alpha-V b)$
$\alpha=V b \Rightarrow x^{t} a_{i} a_{j}^{t} x=b_{i} b_{j} \quad \forall x \in \Omega$
$\diamond$ Penality term
$W=\frac{1}{2} A^{t}$ and $\alpha=-b \Rightarrow F(x)=x^{t} A^{t} A x-2 x^{t} A^{t} b$
$(A x-b)^{2}=0 \forall x \in \Omega$

## Convexifying the objective function of $(\mathrm{Pb})$

Consequence of Debreu's lemma [Lemaréchal, Oustry, 2001]: If $Q$ is positive definite over $\operatorname{Ker}(A)$ then there exists a matrix $V$ such that $Q+A^{t} V A \succcurlyeq 0$ (not true when $Q$ is not definite)

We prove: If $Q$ is positive semidefinite over $\operatorname{Ker}(A)$, there exists $W$ such that $Q+A^{t} W^{t}+W A \succcurlyeq 0$ over $\Re^{n}$

Theorem. Let $A$ and $Q$ be respectively a $p \times n$ matrix and a $n \times n$ symmetric matrix. If $Q$ is positive semidefinite over $\operatorname{Ker}(A)$ then there exists a linear combination of $q_{i j}(x)=x_{i}\left(a_{j}^{t} x-b_{j}\right)$ for $i \in\{1, \ldots, n\}$ and $j \in\{1, \ldots, p\}$ that convexifies the quadratic form $x^{t} Q x$ over $\Re^{n}$. (null functions over $\Omega=\left\{x \in \Re^{n}: A x=b\right\}$ )

Remark: such a combination can be obtained in $O\left(p n^{2}\right)$ time (constructive proof)

Sketch proof (induction on $p$ )
Choose $H=U^{t} W^{t}+W U$ with $U U^{t}=I$ and $U=S^{t} A$. $H$ decomposes into the sum $H=\sum_{i=1}^{p} H_{i}$, where $H_{i}=u_{i} \omega_{i}^{t}+\omega_{i} u_{i}^{t}$. Let $B$ be a matrix whose columns are the vectors of an orthonormal basis of $\operatorname{Ker}(A)$.

$$
L_{i}=\left\{y=\sum_{j=i+1}^{p} z_{j} u_{j}+B z: z_{j} \in \Re \forall j \in\{i+1, \ldots, p\}, z \in \Re^{n-p}\right\}
$$

In particular $L_{p}=\operatorname{Ker}(A)$ and $L_{0}=\Re^{n}$.
Lemma. Let $1 \leq i \leq p$ and $Q$ a $n \times n$ symmetric matrix, if $Q+\sum_{j=i+1}^{p} H_{j} \succcurlyeq 0$ over $L_{i}$ then there exists $\omega_{i}$ such that $Q+\sum_{j=i+1}^{p} H_{j}+u_{i} \omega_{i}^{t}+\omega_{i} u_{i}^{t} \succcurlyeq 0$ over $L_{i-1}$.

## Using convexification in the Lagrangian Approach

$$
\text { (P) } \min _{x} x^{t} Q x+c^{t} x \text { s.t. } x \in \Omega=\{x: A x=b\}
$$

Lemma. If $(P)$ has a solution and $F(x)$ (a constant quadratic function over $\Omega$ ) convexifies the objective function of $(P)$ then there exists $\lambda$ such that
$\min _{x \text { s.t. } A x=b} x^{t} Q x+c^{t} x+F(x)$
$=\min _{x} x^{t} Q x+c^{t} x+F(x)+\lambda^{t}(A x-b)$.

This convexification process transforms the constrained problem $(P)$ into an unconstrained one

## (DP): Partial Lagrangian dual problem of (Pb)

(DP)
$\max _{\mu} \min _{x \text { s.t. } A x=b} x^{t}\left(Q+\sum_{i \in I} \mu_{i} B_{i}\right) x+\left(c+\sum_{i \in I} \mu_{i} d_{i}\right)^{t} x-\sum_{i \in I} \mu_{i} e_{i}$

Partial Lagrangian dual problem of (Pb) where the linear equality constraints are not relaxed.
$\mathfrak{J}=\left\{f_{j}(x)=x^{t} C_{j} x+q_{j}^{t} x+\alpha_{j}: j \in J\right\}:$
A set of null quadratic functions over $\Omega$
$(\mathrm{Pb})_{\mathfrak{J}}:$ add the redundant constraints $f_{j}(x)=0$ to $(\mathrm{Pb})$.
$(D P)_{\mathfrak{J}}$ : partial Lagrangian dual of $(\mathrm{Pb})_{\mathfrak{J}}$ where the constraints $A x=b$ are not relaxed.
$(\mathrm{DP})_{\mathfrak{J}} \Leftrightarrow(\mathrm{DP})$ because $f_{j}(x)=0$ for $x \in \Omega$ and $\forall j \in J$.
(DT) $)_{\mathfrak{J}}$ : total Lagrangian dual of $(P b)_{\mathfrak{J}}$ where all the constraints are relaxed:

$$
\max _{\mu, \nu, \lambda} \min _{x} x^{t} Q(\mu) x+c^{t}(\mu) x-e(\mu)+\sum_{j \in J} \nu_{j} f_{j}(x)+\lambda^{t}(A x-b)
$$

$(\mathrm{DT})_{\mathfrak{J}} \leq(\mathrm{DP})_{\mathfrak{J}}$
Lemma. (convexification) Let $\mu^{*}$ be a solution of (DP). If there exists $\nu^{*}$ such that $\sum_{j \in J} \nu_{j}^{*} f_{j}(x)$ convexifies $x^{t} Q\left(\mu^{*}\right) x+$ $c^{t}\left(\mu^{*}\right) x$, then $(D T)_{\mathfrak{J}} \Leftrightarrow(D P)$.

Recall that we have: $\mathfrak{P}=\left\{x_{i} a_{j}^{t} x-b_{j} x_{i}: \forall i \in\{1, \ldots, n\} \forall j \in\{1, \ldots, p\}\right\}$ Null quadratic functions over $\Omega$ which convexify $x^{t}\left(Q+\sum_{i \in I} \mu_{i} B_{i}\right) x$.

Thus the convexification Lemma implies: (DT) $\mathfrak{P} \Leftrightarrow$ (DP)
Remark. This is not the case with: $\mathfrak{C}=\left\{(A x-b)^{t}(A x-b)\right\}$

## Semidefinite formulation of (DT) $\mathcal{J}_{\mathfrak{J}}$

The semidefinite dual of $(D T)_{\mathfrak{J}}$ is [Lemaréchal, Oustry 2001]:
$(S D P)_{\mathfrak{J}} \quad \min _{X \succcurlyeq x x^{t}} Q \bullet X+c^{t} x$ s.t. $\begin{cases}B_{i} \bullet X+d_{i}^{t} x=e_{i} & i \in I= \\ B_{i} \bullet X+d_{i}^{t} x \leq e_{i} & i \in I \leq \\ A x=b & \\ C_{j} \bullet X+q_{j}^{t} x+\alpha_{j}=0 & j \in J\end{cases}$
Thus $(\mathrm{DT})_{\mathfrak{P}} \Leftrightarrow(\mathrm{DP}) \Leftrightarrow(\mathrm{SDP})_{\mathfrak{P}}$

$$
\mathfrak{C}=\left\{(A x-b)^{t}(A x-b)\right\} \mathfrak{P}=\left\{x_{i} a_{j}^{t} x-b_{j} x_{i}: \forall i \in\{1, \ldots, n\} \forall j \in\{1, \ldots, p\}\right\}
$$

Proposition. $(S D P)_{\mathfrak{C}}$ and $(S D P)_{\mathfrak{F}}$ are equivalent.
$(S D P)_{\mathfrak{C}} \min _{X \succcurlyeq x x^{t}} Q \bullet X+c^{t} x$ s.t. $\begin{cases}B_{i} \bullet X+d_{i}^{t} x=e_{i} & i \in I^{=} \\ B_{i} \bullet X+d_{i}^{t} x \leq e_{i} & i \in I \leq \\ (A x=b) \\ A^{t} A \bullet X-2 b^{t} A x+b^{2}=0\end{cases}$
$(S D P)_{\mathfrak{F}} \quad \min _{X \succcurlyeq x x^{t}} Q \bullet X+c^{t} x$ s.t. $\begin{cases}B_{i} \bullet X+d_{i}^{t} x=e_{i} & i \in I= \\ B_{i} \bullet X+d_{i}^{t} x \leq e_{i} & i \in I \leq \\ A x=b \\ \sum_{k=1}^{n} A_{j k} X_{k i}-b_{j} x_{i}=0 & i \in\{1, \ldots, n\} \\ & j \in\{1, \ldots, p\}\end{cases}$

## Sketch proof

$(X, x)$ feasible of $(S D P)_{\mathfrak{P}}$.
For each $j$ multiply $\sum_{k} A_{j k} X_{k i}-b_{j} x_{i}=0$ by $A_{j i}$, then sum up them all over $j$ and $i$ : $A^{t} A \bullet X-2 b^{t} A x+b^{2}=0$
$(X, x)$ feasible for $(S D P)_{\mathfrak{C}}$.
$A^{t} A \bullet\left(X-x x^{t}\right)+(A x-b)^{2}=0$.
$A^{t} A \bullet\left(X-x x^{t}\right)=0 \Rightarrow A^{t} A\left(X-x x^{t}\right)=0$.
$\forall r, i \in\{1, \ldots, n\} \sum_{k=1}^{n} \sum_{j=1}^{p} A_{j r} A_{j k} X_{k i}=x_{i} \sum_{k=1}^{n} \sum_{j=1}^{p} A_{j r} A_{j k} x_{k}$
$\Rightarrow \forall r, i \in\{1, \ldots, n\} \sum_{j=1}^{p} A_{j r}\left(\sum_{k=1}^{n} A_{j k} X_{k i}-b_{j} x_{i}\right)=0$.
This a linear combination of the $p$ rows of $A$.
$\operatorname{rank}(A)=p \Rightarrow \sum_{k=1}^{n} A_{j k} X_{k i}-b_{j} x_{i}=0 \forall j$ and $\forall i$.

$$
\begin{aligned}
& \mathfrak{P}=\left\{x_{i} a_{j}^{t} x-b_{j} x_{i}: \forall i \in\{1, \ldots, n\} \forall j \in\{1, \ldots, p\}\right\} \\
& \mathfrak{C}=\left\{(A x-b)^{t}(A x-b)\right\}
\end{aligned}
$$


$\longleftrightarrow$ Equivalent programs
$\longrightarrow$ Less or equal to
$<>$ Equivalent if superconsistant

## Conclusion

$\diamond$ A complete characterization of constant quadratic functions over $\left\{x \in \Re^{n}: A x=b\right\}$.
$\diamond$ For general quadratic programs we have $(S D P)_{\mathfrak{P}}=(D T)_{\mathfrak{P}}=(D P)$.
$\diamond$ For Boolean problems (DT) $)_{\mathfrak{C}}=(D P)$ but the supremum is not always reached: consequences for some SDP solvers.
$\diamond$ Better design of semidefinite relaxations of quadratic programs.
$\diamond$ Among the $p \times n$ constraints of (SDP) $\mathfrak{P}_{\mathfrak{P}}$, some may be not active, and thus it would be interesting to foresee which constraints are useful for a given problem.

