Partial Lagrangian relaxation for General Quadratic Programming

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SCRO/JOPT May 9, 2006

Main results (to appear in 4OR)

- 1. A complete characterization of constant quadratic functions over an affine variety $\Omega = \{x \in \Re^n : Ax = b\}.$
- 2. How to convexify the objective function of a general quadratic programming problem (Pb) by using the linear constraints.
- 3. Formulation as a semidefinite program of the partial Lagrangian relaxation of (Pb) where the linear constraints are not relaxed.
- 4. Comparison of two semidefinite relaxations made from two sets of null quadratic functions over an affine variety.

General Quadratic Programming

(Pb) min
$$f(x) = x^t Q x + c^t x$$
 s.t.
$$\begin{cases} x^t B_i x + d_i^t x = e_i & i \in I^= \\ x^t B_i x + d_i^t x \le e_i & i \in I^\le \\ Ax = b \end{cases}$$

 \diamond Boolean quadratic problems can be formulated as (Pb) by considering $x_i^2 = x_i$ for all i in $\{1,\ldots,n\}$

 ♦ Standard semidefinite relaxations of 0-1 quadratic programs are nothing but particular instances of Lagrangian duality [Poljak et al,1995] [Lemaréchal,2003]

♦ For the Boolean case, the *supremum* of the corresponding augmented Lagrangian function is equal to the value of the partial Lagrangian dual (the linear constraints are not relaxed)
 [Lemaréchal, Oustry, 2001]

Characterization of constant quadratic functions over $\Omega = \{x \in \Re^n : Ax = b\}$

Let $F(x) = x^t H x + g^t x$ be a quadratic function that takes a constant value over Ω . (underlying idea: add redundant constraints)

For any $x \in \Omega$, $u \in Ker(A) = \{u \in \Re^n : Au = 0\}$, and $\lambda \in \Re$ $F(x + \lambda u) = F(x) = F(x) + \lambda^2 u^t Hu + 2\lambda u^t Hx + \lambda g^t u$ i.e. $F(x + \lambda u)$ does not depend on λ .

Necessary and sufficient conditions on F(x):

$$u^{t}Hu = 0 \ \forall u \in Ker(A)$$

$$[A]$$

$$2u^{t}Hx + g^{t}u = 0 \ \forall u \in Ker(A) \ \forall x \in \Omega$$

$$[B]$$

Characterization of constant quadratic functions over $\Omega = \{x \in \Re^n : Ax = b\}$

Lemma. $q(u) = u^t H u$ is a null quadratic form over Ker(A) if and only if $q(u) = u^t (A^t W^t + WA) u$, where W is any $n \times p$ -matrix.

Sketch proof. choose a "good" basis for the quadratic form, write and simplify P^tHP , then obtain $H = (P^t)^{-1}P^tHPP^{-1}$. $P = \begin{bmatrix} A^t & B \end{bmatrix}$, the n - p columns of B are a basis of Ker(A).

Theorem. $F(x) = x^t Hx + g^t x$ is a constant quadratic function over $\{x : Ax = b\}$ if and only if $F(x) = x^t (A^t W^t + WA) x + (A^t \alpha - 2Wb)^t x$, where W is a $n \times p$ -matrix and α is a p-vector.

Any constant quadratic function over $\Omega = \{x \in \Re^n : Ax = b\}$ can be obtained by setting particular values to W and α in

$$F(x) = x^t \left(A^t W^t + W A \right) x + \left(A^t \alpha - 2W b \right)^t x$$

◇ Product member by variable $W = \frac{1}{2}E^{ij}$ and α = 0 where $E^{ij}_{ij} = 1$ otherwise $E^{kl}_{ij} = 0$. $F(x) = x_i a^t_j x - b_j x_i = 0$ ∀x ∈ Ω

♦ Product member by member $W = \frac{1}{2}A^tV$ with $V = \frac{1}{2}(E^{ij} + E^{ji})$ $F(x) = x^t a_i a_j^t x + x^t A^t (\alpha - Vb)$ $\alpha = Vb \Rightarrow x^t a_i a_j^t x = b_i b_j$ ∀x ∈ Ω

♦ Penality term

$$W = \frac{1}{2}A^t$$
 and $\alpha = -b \Rightarrow F(x) = x^t A^t A x - 2x^t A^t b$
 $(Ax - b)^2 = 0 \forall x \in Ω$

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Convexifying the objective function of (Pb)

Consequence of Debreu's lemma [Lemaréchal, Oustry, 2001]: If Q is positive definite over Ker(A) then there exists a matrix V such that $Q + A^t VA \succeq 0$ (not true when Q is not definite)

We prove: If Q is positive semidefinite over Ker(A), there exists W such that $Q + A^tW^t + WA \succeq 0$ over \Re^n

Theorem. Let A and Q be respectively a $p \times n$ matrix and a $n \times n$ symmetric matrix. If Q is positive semidefinite over Ker(A) then there exists a linear combination of $q_{ij}(x) = x_i(a_j^t x - b_j)$ for $i \in \{1, ..., n\}$ and $j \in \{1, ..., p\}$ that convexifies the quadratic form $x^t Qx$ over \Re^n . (null functions over $\Omega = \{x \in \Re^n : Ax = b\}$)

Remark: such a combination can be obtained in $O(pn^2)$ time (constructive proof)

Sketch proof (induction on p)

Choose $H = U^t W^t + WU$ with $UU^t = I$ and $U = S^t A$. H decomposes into the sum $H = \sum_{i=1}^p H_i$, where $H_i = u_i \omega_i^t + \omega_i u_i^t$. Let B be a matrix whose columns are the vectors of an orthonormal basis of Ker(A).

$$L_{i} = \left\{ y = \sum_{j=i+1}^{p} z_{j}u_{j} + Bz : z_{j} \in \Re \ \forall j \in \{i+1,...,p\}, z \in \Re^{n-p} \right\}$$

In particular $L_p = Ker(A)$ and $L_0 = \Re^n$.

Lemma. Let $1 \leq i \leq p$ and Q a $n \times n$ symmetric matrix, if $Q + \sum_{j=i+1}^{p} H_j \geq 0$ over L_i then there exists ω_i such that $Q + \sum_{j=i+1}^{p} H_j + u_i \omega_i^t + \omega_i u_i^t \geq 0$ over L_{i-1} .

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Using convexification in the Lagrangian Approach

(P)
$$\min_{x} x^{t}Qx + c^{t}x$$
 s.t. $x \in \Omega = \{x : Ax = b\}$

Lemma. If (P) has a solution and F(x) (a constant quadratic function over Ω) convexifies the objective function of (P) then there exists λ such that $\min_{x \text{ s.t. } Ax = b} x^t Qx + c^t x + F(x)$ $= \min_x x^t Qx + c^t x + F(x) + \lambda^t (Ax - b).$

This convexification process transforms the *constrained* problem (P) into an *unconstrained* one

(DP): Partial Lagrangian dual problem of (Pb)

(DP) $\max_{\mu} \min_{x \text{ s.t. } Ax=b} x^{t} \left(Q + \sum_{i \in I} \mu_{i} B_{i}\right) x + \left(c + \sum_{i \in I} \mu_{i} d_{i}\right)^{t} x - \sum_{i \in I} \mu_{i} e_{i}$

Partial Lagrangian dual problem of (Pb) where the linear equality constraints are not relaxed.

 $\mathfrak{J} = \left\{ f_j(x) = x^t C_j x + q_j^t x + \alpha_j : j \in J \right\}:$ A set of null quadratic functions over Ω

 $(Pb)_{\mathfrak{J}}$: add the redundant constraints $f_j(x) = 0$ to (Pb).

 $(\mathsf{DP})_{\mathfrak{J}}$: partial Lagrangian dual of $(\mathsf{Pb})_{\mathfrak{J}}$ where the constraints Ax = b are not relaxed. $(\mathsf{DP})_{\mathfrak{J}} \Leftrightarrow (\mathsf{DP})$ because $f_j(x) = 0$ for $x \in \Omega$ and $\forall j \in J$.

 $(DT)_{\mathfrak{J}}$: total Lagrangian dual of $(Pb)_{\mathfrak{J}}$ where all the constraints are relaxed:

$$\max_{\mu,\nu,\lambda} \min_{x} x^{t} Q(\mu) x + c^{t}(\mu) x - e(\mu) + \sum_{j \in J} \nu_{j} f_{j}(x) + \lambda^{t} (Ax - b)$$
$$(\mathsf{DT})_{\mathfrak{J}} \leq (\mathsf{DP})_{\mathfrak{J}}$$

Lemma. (convexification) Let μ^* be a solution of (DP). If there exists ν^* such that $\sum_{j \in J} \nu_j^* f_j(x)$ convexifies $x^t Q(\mu^*) x + c^t(\mu^*) x$, then $(DT)_{\mathfrak{J}} \Leftrightarrow (DP)$. **Recall that we have**: $\mathfrak{P} = \left\{ x_i a_j^t x - b_j x_i : \forall i \in \{1, ..., n\} \ \forall j \in \{1, ..., p\} \right\}$ Null quadratic functions over Ω which convexify $x^t (Q + \sum_{i \in I} \mu_i B_i) x$.

Thus the convexification Lemma implies: $(DT)_{\mathfrak{P}} \Leftrightarrow (DP)$

Remark. This is not the case with: $\mathfrak{C} = \{ (Ax - b)^t (Ax - b) \}$

Semidefinite formulation of $(DT)_{\tilde{J}}$

The semidefinite dual of $(DT)_{\mathfrak{J}}$ is [Lemaréchal, Oustry 2001]:

$$(SDP)_{\mathfrak{J}} \qquad \min_{X \succcurlyeq xx^{t}} Q \bullet X + c^{t}x \text{ s.t.} \begin{cases} B_{i} \bullet X + d_{i}^{t}x = e_{i} & i \in I^{=} \\ B_{i} \bullet X + d_{i}^{t}x \leq e_{i} & i \in I^{\leq} \\ Ax = b \\ C_{j} \bullet X + q_{j}^{t}x + \alpha_{j} = 0 & j \in J \end{cases}$$

Thus $(DT)_{\mathfrak{P}} \Leftrightarrow (DP) \Leftrightarrow (SDP)_{\mathfrak{P}}$

$$\mathfrak{C} = \left\{ (Ax - b)^t (Ax - b) \right\} \mathfrak{P} = \left\{ x_i a_j^t x - b_j x_i : \forall i \in \{1, ..., n\} \ \forall j \in \{1, ..., p\} \right\}$$

Proposition. $(SDP)_{\mathfrak{C}}$ and $(SDP)_{\mathfrak{P}}$ are equivalent.

$$(SDP)_{\mathfrak{C}} \qquad \min_{X \succcurlyeq xx^{t}} Q \bullet X + c^{t}x \text{ s.t.} \begin{cases} B_{i} \bullet X + d_{i}^{t}x = e_{i} & i \in I^{=} \\ B_{i} \bullet X + d_{i}^{t}x \leq e_{i} & i \in I^{\leq} \\ (Ax = b) \\ A^{t}A \bullet X - 2b^{t}Ax + b^{2} = 0 \end{cases}$$
$$(SDP)_{\mathfrak{P}} \qquad \min_{X \succcurlyeq xx^{t}} Q \bullet X + c^{t}x \text{ s.t.} \begin{cases} B_{i} \bullet X + d_{i}^{t}x = e_{i} & i \in I^{=} \\ B_{i} \bullet X + d_{i}^{t}x \leq e_{i} & i \in I^{\leq} \\ Ax = b \\ \sum_{k=1}^{n} A_{jk}X_{ki} - b_{j}x_{i} = 0 & i \in \{1, ..., n\} \\ j \in \{1, ..., p\} \end{cases}$$

Sketch proof

(X, x) feasible of $(SDP)_{\mathfrak{P}}$. For each j multiply $\sum_k A_{jk}X_{ki} - b_jx_i = 0$ by A_{ji} , then sum up them all over j and i: $A^tA \bullet X - 2b^tAx + b^2 = 0$

$$(X, x) \text{ feasible for } (SDP)_{\mathfrak{C}}.$$

$$A^{t}A \bullet (X - xx^{t}) + (Ax - b)^{2} = 0.$$

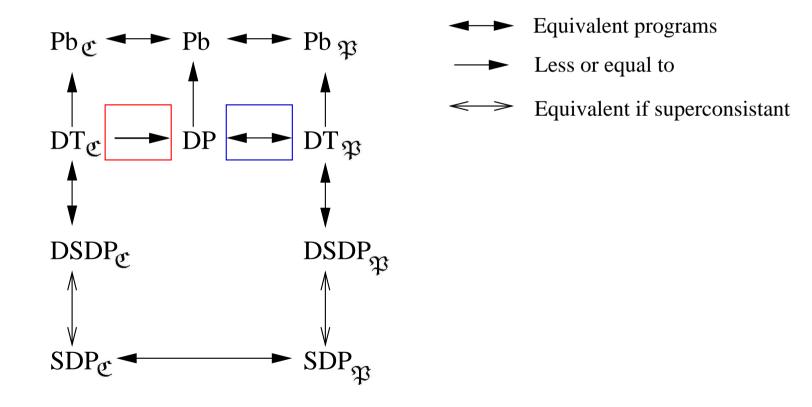
$$A^{t}A \bullet (X - xx^{t}) = 0 \Rightarrow A^{t}A (X - xx^{t}) = 0.$$

$$\forall r, i \in \{1, ..., n\} \sum_{k=1}^{n} \sum_{j=1}^{p} A_{jr}A_{jk}X_{ki} = x_{i} \sum_{k=1}^{n} \sum_{j=1}^{p} A_{jr}A_{jk}x_{k}$$

$$\Rightarrow \forall r, i \in \{1, ..., n\} \sum_{j=1}^{p} A_{jr} (\sum_{k=1}^{n} A_{jk}X_{ki} - b_{j}x_{i}) = 0.$$
This a linear combination of the *p* rows of *A*.

$$rank(A) = p \Rightarrow \sum_{k=1}^{n} A_{jk}X_{ki} - b_{j}x_{i} = 0 \ \forall j \text{ and } \forall i.$$

$$\mathfrak{P} = \left\{ x_i a_j^t x - b_j x_i : \forall i \in \{1, ..., n\} \ \forall j \in \{1, ..., p\} \right\}$$
$$\mathfrak{C} = \left\{ (Ax - b)^t (Ax - b) \right\}$$



Conclusion

 \diamond A complete characterization of constant quadratic functions over $\{x \in \Re^n : Ax = b\}$.

 \Diamond For general quadratic programs we have $(SDP)_{\mathfrak{P}} = (DT)_{\mathfrak{P}} = (DP)$.

 \diamond For Boolean problems $(DT)_{\mathfrak{C}} = (DP)$ but the supremum is not always reached: consequences for some SDP solvers.

♦ Better design of semidefinite relaxations of quadratic programs.

 \diamond Among the $p \times n$ constraints of $(SDP)_{\mathfrak{P}}$, some may be not active, and thus it would be interesting to foresee which constraints are useful for a given problem.